

**Proposed Update** Unicode Standard Annex #9**UNICODE BIDIRECTIONAL ALGORITHM**

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Summary

This annex describes specifications for the positioning of characters flowing from right to left, such as Arabic or Hebrew.

Status

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1 Introduction

The Unicode Standard prescribes a *memory* representation order known as logical order. When text is presented in horizontal lines, most scripts display characters from left to right. However, there are several scripts (such as Arabic or Hebrew) where the natural ordering of horizontal text in display is from right to left. If all of the text has the same horizontal direction, then the ordering of the display text is unambiguous. However, when bidirectional text (a mixture of left-to-right and right-to-left horizontal text) is present, some ambiguities can arise in determining the ordering of the displayed characters.

This annex describes the algorithm used to determine the directionality for bidirectional Unicode text. The algorithm extends the implicit model currently employed by a number of existing implementations and adds explicit format codes for special circumstances. In most cases, there is no need to include additional information with the text to obtain correct display ordering.

However, in the case of bidirectional text, there are circumstances where an implicit bidirectional ordering is not sufficient to produce comprehensible text. To deal with these cases, a minimal set of directional formatting codes is defined to control the ordering of characters when rendered. This allows exact control of the display ordering for legible interchange and ensures that plain text used for simple items like filenames or labels can always be correctly ordered for display.

The directional formatting codes are used *only* to influence the display ordering of text. In all other respects they should be ignored—they have no effect on the comparison of text or on word breaks, parsing, or numeric analysis.

When working with bidirectional text, the characters are still interpreted in logical order—only the display is affected. The display ordering of bidirectional text depends on the directional properties of the characters in the text. Note that there are important security issues connected with bidirectional text: for more information, see [UTR36](#).

2 Directional Formatting Codes

Two types of explicit codes are used to modify the standard implicit Unicode Bidirectional Algorithm. In addition, there are implicit ordering codes, the *right-to-left* and *left-to-right* marks. All of these codes are limited to the current paragraph; thus their effects are terminated by a *paragraph separator*. The directional types left-to-right and right-to-left are called *strong types*, and characters of those types are called strong directional characters. The directional types associated with numbers are called *weak types*, and characters of those types are called weak directional characters.



These controls all have the property *Bidi_Control*.

U+200E..U+200F LEFT-TO-RIGHT MARK..RIGHT-TO-LEFT MARK
U+202A..U+202E LEFT-TO-RIGHT EMBEDDING..RIGHT-TO-LEFT OVERRIDE

Although the term *embedding* is used for some explicit codes, the text within the scope of the codes is not independent of the surrounding text. Characters within an embedding can affect the ordering of characters outside, and vice versa. The algorithm is designed so that the use of explicit codes can be equivalently represented by out-of-line information, such as stylesheet information. However, any alternative representation will be defined by reference to the behavior of the explicit codes in this algorithm.

2.1 Explicit Directional Embedding



The following codes signal that a piece of text is to be treated as embedded. For example, an English quotation in the middle of an Arabic sentence could be marked as being embedded left-to-right text. If there were a Hebrew phrase in the middle of the English quotation, the that phrase could be marked as being embedded right-to-left. These codes allow for nested embeddings.

Abbr.	Code	Chart	Name	Description
LRE	U+202A		LEFT-TO-RIGHT EMBEDDING	Treat the following text as embedded left-to-right.
RLE	U+202B		RIGHT-TO-LEFT EMBEDDING	Treat the following text as embedded right-to-left.

The effect of right-left line direction, for example, can be accomplished by embedding the text with RLE...PDF.
On web pages, these characters should be replaced by using the *dir* attribute with the values *dir="ltr"* or *dir="rtl"*. For more information, see [UTR20](#).

2.2 Explicit Directional Overrides


The following codes allow the bidirectional character types to be overridden when required for special cases, such as for part numbers. These codes allow for nested directional overrides.
These characters are to be avoided wherever possible, because of security concerns. For more information, see [UTR36](#).

Abbr.	Code	Chart	Name	Description
LRO	U+202D		LEFT-TO-RIGHT OVERRIDE	Force following characters to be treated as strong left-to-right characters.
RLO	U+202E		RIGHT-TO-LEFT OVERRIDE	Force following characters to be treated as strong right-to-left characters.

The precise meaning of these codes will be made clear in the discussion of the algorithm. The right-to-left override, for example, can be used to force a part number made of mixed English, digits and Hebrew letters to be written from right to left.

2.3 Terminating Explicit Directional Code



The following code terminates the effects of the last explicit code (either embedding or override) and restores the bidirectional state to what it was before that code was encountered.

Abbr.	Code	Chart	Name	Description
PDF	U+202C		POP DIRECTIONAL FORMATTING	Restore the bidirectional state to what it was before the last LRE, RLE, RLO, or LRO.

The precise meaning of this code will be made clear in the discussion of the algorithm.

2.4 Implicit Directional Marks

These characters are very light-weight codes. They act exactly like right-to-left or left-to-right characters, except that they do not display or have any other semantic effect. Their use is more convenient than using explicit embeddings or overrides because their scope is much more local.

Abbr.	Code	Chart	Name	Description
LRM	U+200E		LEFT-TO-RIGHT MARK	Left-to-right zero-width character
RLM	U+200F		RIGHT-TO-LEFT MARK	Right-to-left zero-width character

There is no special mention of the implicit directional marks in the following algorithm. That is because their effect on bidirectional ordering is exactly the same as a corresponding strong directional character; the only difference is that they do not appear in the display.

3 Basic Display Algorithm

The Bidirectional Algorithm takes a stream of text as input and proceeds in three main phases:

- Separation of the input text into paragraphs. The rest of the algorithm affects only the text between paragraph separators.
- Resolution of the embedding levels of the text. In this phase, the directional character types, plus the explicit format codes, are used to produce resolved embedding levels.
- Reordering the text for display on a line-by-line basis using the resolved embedding levels, once the text has been broken into lines.

The algorithm reorders text only within a paragraph; characters in one paragraph have no effect on characters in a different paragraph. Paragraphs are divided by the Paragraph Separator or appropriate Newline Function (for guidelines on the handling of CR, LF, and CRLF, see *Section 4.4, Directionality*, and *Section 5.8, Newline Guidelines* of [\[Unicode\]](#)). Paragraphs may also be determined by higher-level protocols: for example, the text in two different cells of a table will be in different paragraphs.

Combining characters always attach to the preceding base character in the memory representation. Even after reordering for display and performing character shaping, the glyph representing a combining character will attach to the glyph representing its base character in memory. Depending on the line orientation and the placement direction of base letterform glyphs, it may, for example, attach to the glyph on the left, or on the right, or above.

This annex uses the numbering conventions for normative definitions and rules in *Table 1*.

Table 1. Normative Definitions and Rules

Numbering	Section
BDn	Definitions
Pn	Paragraph levels
Xn	Explicit levels and directions
Wn	Weak types
Nn	Neutral types
In	Implicit levels
Ln	Resolved levels

3.1 Definitions

BD1. The *bidirectional characters types* are values assigned to each Unicode character, including unassigned characters.

BD2. *Embedding levels*

are numbers that indicate how deeply the text is nested, and the default direction of text on that level. The minimum embedding level of text is zero, and the maximum explicit depth is level 61.

Embedding levels are explicitly set by both override format codes and by embedding format codes; higher numbers mean the text is more deeply nested. The reason for having a limitation is to provide a precise stack limit for implementations to guarantee the same results. Sixty-one levels is far more than sufficient for ordering, even with mechanically generated formatting; the display becomes rather muddled with more than a small number of embeddings.

BD3. The default direction of the current embedding level (for the character in question) is called the *embedding direction*. It is **L** if the embedding level is even, and **R** if the embedding level is odd.

For example, in a particular piece of text, Level 0 is plain English text. Level 1 is plain Arabic text, possibly embedded within English level 0 text. Level 2 is English text, possibly embedded within Arabic level 1 text, and so on. Unless their direction is overridden, English text and numbers will always be an even level; Arabic text (excluding numbers) will always be an odd level. The exact meaning of the embedding level will become clear when the reordering algorithm is discussed, but the following

provides an example of how the algorithm works.

BD4. The *paragraph embedding level* is the embedding level that determines the default bidirectional orientation of the text in that paragraph.

BD5. The direction of the paragraph embedding level is called the *paragraph direction*.

- In some contexts the paragraph direction is also known as the *base direction*.

BD6. The *directional override status* determines whether the bidirectional type of characters is to be reset. The override status is set by using explicit directional controls. This status has three states, as shown in *Table 2*.

Table 2. Directional Override Status

Status	Interpretation
Neutral	No override is currently active
Right-to-left	Characters are to be reset to R
Left-to-right	Characters are to be reset to L

BD7. A *level run*

is a maximal substring of characters that have the same embedding level. It is maximal in that no character immediately before or after the substring has the same level (a level run is also known as a *directional run*).

Example

In this and the following examples, case is used to indicate different implicit character types for those unfamiliar with right-to-left letters. Uppercase letters stand for right-to-left characters (such as Arabic or Hebrew), and lowercase letters stand for left-to-right characters (such as English or Russian).

Memory: car is THE CAR in arabic

Character types: LLL-LL-RRR-RRR-LL-LLLLLL

Resolved levels: 000000011111110000000000

Notice that the neutral character (space) between THE and CAR gets the level of the surrounding characters. The level of the neutral characters can also be changed by inserting appropriate directional marks around neutral characters. These marks have no other effects.

Table 3

lists additional abbreviations used in the examples and internal character types used in the algorithm.

Table 3. Abbreviations for Examples and Internal Types

Symbol	Description
N	Neutral or Separator (B, S, WS, ON)
e	The text ordering type (L or R) that matches the embedding level direction (even or odd)
sor	The text ordering type (L or R) assigned to the position before a level run.
eor	The text ordering type (L or R) assigned to the position after a level run.

3.2 Bidirectional Character Types

The normative bidirectional character types for each character are specified in the [Unicode Character Database \[UCD\]](#) and are summarized in *Table 4*. This is a summary only: there are exceptions to the general scope. For example, certain characters such as U+0CBF KANNADA VOWEL SIGN I are given Type L (instead of NSM) to preserve canonical equivalence.

- The term European digits is used to refer to decimal forms common in Europe and elsewhere, and Arabic-Indic digits to refer to the native Arabic forms. (See *Section 8.2, Arabic* of [\[Unicode\]](#), for more details on naming digits.)
- Unassigned characters are given strong types in the algorithm. This is an explicit exception to the general Unicode conformance requirements with respect to unassigned characters. As characters become assigned in the future, these bidirectional types may change. For assignments to character types, see `DerivedBidiClass.txt` [\[DerivedBIDI\]](#) in the [\[UCD\]](#).
- Private-use characters can be assigned different values by a conformant implementation.
- For the purpose of the Bidirectional Algorithm, inline objects (such as graphics) are treated as if they are an U+FFFC OBJECT REPLACEMENT CHARACTER.
- As of Unicode 4.0, the Bidirectional Character Types of a few Indic characters were altered so that the Bidirectional Algorithm preserves canonical equivalence. That is, two canonically equivalent strings will result in equivalent ordering after applying the algorithm. This invariant will be maintained in the future.

Note: The Bidirectional Algorithm does *not* preserve compatibility equivalence.

Table 4. Bidirectional Character Types

Category	Type	Description	General Scope
Strong	L	Left-to-Right	LRM, most alphabetic, syllabic, Han ideographs, non-European or non-Arabic digits, ...
	LRE	Left-to-Right Embedding	LRE
	LRO	Left-to-Right Override	LRO
	R	Right-to-Left	RLM, Hebrew alphabet, and related punctuation
	AL	Right-to-Left Arabic	Arabic, Thaana, and Syriac alphabets, most punctuation specific to those scripts, ...
	RLE	Right-to-Left Embedding	RLE
	RLO	Right-to-Left Override	RLO
Weak	PDF	Pop Directional Format	PDF
	EN	European Number	European digits, Eastern Arabic-Indic digits, ...
	ES	European Number Separator	PLUS SIGN, MINUS SIGN
	ET	European Number Terminator	DEGREE SIGN, currency symbols, ...
	AN	Arabic Number	Arabic-Indic digits, Arabic decimal and thousands separators, ...
	CS	Common Number Separator	COLON, COMMA, FULL STOP (<i>period</i>), NO-BREAK SPACE, ...
	NSM	Nonspacing Mark	Characters marked Mn (Nonspacing_Mark) and Me (Enclosing_Mark) in the Unicode Character Database
	BN	Boundary Neutral	Most formatting and control characters, other than those explicitly given types above
Neutral	B	Paragraph Separator	PARAGRAPH SEPARATOR, appropriate Newline Functions, higher-level protocol paragraph determination
	S	Segment Separator	<i>Tab</i>
	WS	Whitespace	SPACE, FIGURE SPACE, LINE SEPARATOR, FORM FEED, General Punctuation spaces, ...
	ON	Other Neutrals	All other characters, including OBJECT REPLACEMENT CHARACTER

3.3 Resolving Embedding Levels

The body of the Bidirectional Algorithm uses character types and explicit codes to produce a list of resolved levels. This resolution process consists of five steps: (1) determining the paragraph level; (2) determining explicit embedding levels and directions; (3) resolving weak types; (4) resolving neutral types; and (5) resolving implicit embedding levels.

3.3.1 The Paragraph Level

P1. Split the text into separate paragraphs. A paragraph separator is kept with the previous paragraph. Within each paragraph, apply all the other rules of this algorithm.

P2. In each paragraph, find the first character of type L, AL, or R.

Because paragraph separators delimit text in this algorithm, this will generally be the first strong character after a paragraph separator or at the very beginning of the text. Note that the characters of type LRE, LRO, RLE, or RLO are ignored in this rule. This is because typically they are used to indicate that the embedded text is the *opposite* direction than the paragraph level.

P3. If a character is found in P2 and it is of type AL or R, then set the paragraph embedding level to one; otherwise, set it to zero.

Whenever a higher-level protocol specifies the paragraph level, rules P2 and P3 do not apply.

3.3.2 Explicit Levels and Directions

All explicit embedding levels are determined from the embedding and override codes, by applying the explicit level rules X1 through X9. These rules are applied as part of the same logical pass over the input.

Explicit Embeddings

X1. Begin by setting the current embedding level to the paragraph embedding level. Set the directional override status to neutral. Process each character iteratively, applying rules X2 through X9. Only embedding levels from 0 to 61 are valid in this phase.

In the resolution of levels in rules I1 and I2, the maximum embedding level of 62 can be reached.

*X2. With each RLE, compute the least greater **odd** embedding level.*

*a. If this new level would be valid, then this embedding code is valid. Remember (push) the current embedding level and override status. Reset the current level to this new level, and reset the override status to **neutral**.*

b. If the new level would not be valid, then this code is invalid. Do not change the current level or override status.

For example, level 0 → 1; levels 1, 2 → 3; levels 3, 4 → 5; ...59, 60 → 61; above 60, no change (do not change levels with RLE if the new level would be invalid).

*X3. With each LRE, compute the least greater **even** embedding level.*

*a. If this new level would be valid, then this embedding code is valid. Remember (push) the current embedding level and override status. Reset the current level to this new level, and reset the override status to **neutral**.*

b. If the new level would not be valid, then this code is invalid. Do not change the current level or override status.

For example, levels 0, 1 → 2; levels 2, 3 → 4; levels 4, 5 → 6; ...58, 59 → 60; above 59, no

change (do not change levels with LRE if the new level would be invalid).

Explicit Overrides

An explicit directional override sets the embedding level in the same way the explicit embedding codes do, but also changes the directional character type of affected characters to the override direction.

*X4. With each RLO, compute the least greater **odd** embedding level.*

- a. If this new level would be valid, then this embedding code is valid. Remember (push) the current embedding level and override status. Reset the current level to this new level, and reset the override status to **right-to-left**.*
- b. If the new level would not be valid, then this code is invalid. Do not change the current level or override status.*

*X5. With each LRO, compute the least greater **even** embedding level.*

- a. If this new level would be valid, then this embedding code is valid. Remember (push) the current embedding level and override status. Reset the current level to this new level, and reset the override status to **left-to-right**.*
- b. If the new level would not be valid, then this code is invalid. Do not change the current level or override status.*

X6. For all types besides RLE, LRE, RLO, LRO, and PDF:

- a. Set the level of the current character to the current embedding level.*
- b. Whenever the directional override status is not neutral, reset the current character type to the directional override status.*

If the directional override status is neutral, then characters retain their normal types: Arabic characters stay AL, Latin characters stay L, neutrals stay N, and so on. If the directional override status is R, then characters become R. If the directional override status is L, then characters become L.

Terminating Embeddings and Overrides

There is a single code to terminate the scope of the current explicit code, whether an embedding or a directional override. All codes and pushed states are completely popped at the end of paragraphs.

X7. With each PDF, determine the matching embedding or override code. If there was a valid matching code, restore (pop) the last remembered (pushed) embedding level and directional override.

*X8. All explicit directional embeddings and overrides are completely terminated at the end of each paragraph. Paragraph separators are **not** included in the embedding.*

X9. Remove all RLE, LRE, RLO, LRO, PDF, and BN codes.

- Note that an implementation does not have to actually remove the codes; it just has to behave as though the codes were not present for the remainder of the algorithm.

Conformance does not require any particular placement of these codes as long as all other characters are ordered correctly.

See [Section 5, *Implementation Notes*](#), for information on implementing the algorithm without removing the formatting codes.

- The *zero width joiner* and *non-joiner* affect the shaping of the adjacent characters—those that are adjacent in the original backing-store order, even though those characters may end up being rearranged to be non-adjacent by the Bidirectional Algorithm. For more information, see [Section 5.3, *Joiners*](#).

*X10. The remaining rules are applied to each run of characters at the same level. For each run, determine the start-of-level-run (**sor**) and end-of-level-run (**eor**) type, either L or R. This depends on the higher of the two levels on either side of the boundary (at the start or end of the paragraph, the level of the “other” run is the base embedding level). If the higher level is odd, the type is R; otherwise, it is L.*

For example:

```
Levels:  0    0    0    1    1    1    2
Runs:    <--- 1 ---> <--- 2 ---> <3>
```

Run 1 is at level 0, *sor* is L, *eor* is R.

Run 2 is at level 1, *sor* is R, *eor* is L.

Run 3 is at level 2, *sor* is L, *eor* is L.

For two adjacent runs, the *eor* of the first run is the same as the *sor* of the second.

3.3.3 Resolving Weak Types

Weak types are now resolved one level run at a time. At level run boundaries where the type of the character on the other side of the boundary is required, the type assigned to *sor* or *eor* is used.

Nonspacing marks are now resolved based on the previous characters.

*W1. Examine each nonspacing mark (NSM) in the level run, and change the type of the NSM to the type of the previous character. If the NSM is at the start of the level run, it will get the type of **sor**.*

Assume in this example that *sor* is R:

```
AL  NSM NSM → AL  AL  AL
sor NSM      → sor R
```

The text is next parsed for numbers. This pass will change the directional types European Number Separator, European Number Terminator, and Common Number Separator to be European Number text, Arabic Number text, or Other Neutral text. The text to be scanned may have already had its type altered by directional overrides. If so, then it will not parse as numeric.

*W2. Search backward from each instance of a European number until the first strong type (R, L, AL, or **sor**) is found. If an AL is found, change the type of the European number to Arabic number.*

```

AL EN      → AL AN
AL N EN    → AL N AN
sor N EN   → sor N EN
L N EN     → L N EN
R N EN     → R N EN

```

W3. Change all ALs to R.

W4. A single European separator between two European numbers changes to a European number. A single common separator between two numbers of the same type changes to that type.

```

EN ES EN → EN EN EN
EN CS EN → EN EN EN
AN CS AN → AN AN AN

```

W5. A sequence of European terminators adjacent to European numbers changes to all European numbers.

```

ET ET EN → EN EN EN
EN ET ET → EN EN EN
AN ET EN → AN EN EN

```

W6. Otherwise, separators and terminators change to Other Neutral.

```

AN ET      → AN ON
L  ES EN → L  ON EN
EN CS AN → EN ON AN
ET AN      → ON AN

```

*W7. Search backward from each instance of a European number until the first strong type (R, L, or **sor**) is found. If an L is found, then change the type of the European number to L.*

```

L  N EN => L  N  L
R  N EN => R  N  EN

```

3.3.4 Resolving Neutral Types

Neutral types are now resolved one level run at a time. At level run boundaries where the type of the character on the other side of the boundary is required, the type assigned to **sor** or **eor** is used.

The next phase resolves the direction of the neutrals. The results of this phase are that all neutrals become either **R** or **L**. Generally, neutrals take on the direction of the surrounding text. In case of a conflict, they take on the embedding direction.

*N1. A sequence of neutrals takes the direction of the surrounding strong text if the text on both sides has the same direction. European and Arabic numbers act as if they were R in terms of their influence on neutrals. Start-of-level-run (**sor**) and end-of-level-run (**eor**) are*

used at level run boundaries.

```

R  N  R  → R  R  R
L  N  L  → L  L  L
R  N  AN → R  R  AN
AN N  R  → AN R  R
R  N  EN → R  R  EN
EN N  R  → EN R  R

```

Note that any AN or EN remaining after W7 will be in an right-to-left context.

N2. Any remaining neutrals take the embedding direction.

```
N → e
```

Assume in this example that *eor* is L and *sor* is R. Then an application of N1 and N2 yields the following:

```

L  N eor → L  L eor
R  N eor → R  e eor
sor N L   → sor e L
sor N R   → sor R R

```

Examples.

A list of numbers separated by neutrals and embedded in a directional run will come out in the run's order.

Storage: he said "THE VALUES ARE 123, 456, 789, OK".

Display: he said "KO ,789 ,456 ,123 ERA SEULAV EHT".

In this case, both the comma and the space between the numbers take on the direction of the surrounding text (uppercase = right-to-left), ignoring the numbers. The commas are not considered part of the number because they are not surrounded on both sides (see *Section 3.3.3, Resolving Weak Types*). However, if there is an adjacent left-to-right sequence, then European numbers will adopt that direction:

Storage: he said "IT IS A bmw 500, OK."

Display: he said ".KO ,bmw 500 A SI TI".

3.3.5 Resolving Implicit Levels

In the final phase, the embedding level of text may be increased, based on the resolved character type. Right-to-left text will always end up with an odd level, and left-to-right and numeric text will always end up with an even level. In addition, numeric text will always end up with a higher level than the paragraph level. (Note that it is possible for text to end up at levels higher than 61 as a result of this process.) This results in the following rules:

I1. For all characters with an even (left-to-right) embedding direction, those of type R go up one level and those of type AN or EN go up two levels.

I2. For all characters with an odd (right-to-left) embedding direction, those of type L, EN or

AN go up one level.

Table 5 summarizes the results of the implicit algorithm.

Table 5. Resolving Implicit Levels

Type	Embedding Level	
	Even	Odd
L	EL	EL + 1
R	EL + 1	EL
AN	EL + 2	EL + 1
EN	EL + 2	EL + 1

3.4 Reordering Resolved Levels

The following rules describe the logical process of finding the correct display order. As opposed to resolution phases, these rules act on a per-line basis *and are applied **after** any line wrapping is applied to the paragraph.*

Logically there are the following steps:

- The levels of the text are determined according to the previous rules.
- The characters are shaped into glyphs according to their context (*taking the embedding levels into account for mirroring*).
- The accumulated widths of those glyphs (*in logical order*) are used to determine line breaks.
- For each line, rules L1–L4 are used to reorder the characters on that line.
- The glyphs corresponding to the characters on the line are displayed in that order.

L1. On each line, reset the embedding level of the following characters to the paragraph embedding level:

1. *Segment separators,*
 2. *Paragraph separators,*
 3. *Any sequence of whitespace characters preceding a segment separator or paragraph separator, and*
 4. *Any sequence of white space characters at the end of the line.*
- The types of characters used here are the *original* types, not those modified by the previous phase.
 - Because a PARAGRAPH SEPARATOR breaks lines, there will be at most one per line, at the end of that line.

In combination with the following rule, this means that trailing whitespace will appear at the visual end of the line (in the paragraph direction). Tabulation will always have a consistent direction within a paragraph.

L2. From the highest level found in the text to the lowest odd level on each line, including intermediate levels not actually present in the text, reverse any contiguous sequence of characters that are at that level or higher.

This rule reverses a progressively larger series of substrings.



The following examples illustrate the reordering, showing the successive steps in application of Rule [L2](#). The original text, including any embedding codes for producing the particular levels, is shown in the "Storage" row in the example tables. The application of the rules from [Section 3.3 Resolving Embedding Levels](#) and of the Rule [L1](#) results in (a) text with Bidi Controls and BN characters removed, plus (b) resolved levels. These are listed in the rows "Before Reordering" and "Resolved Levels". Each successive row thereafter shows the one pass of reversal from Rule L2, such as "Reverse levels 1-2". At each iteration, the underlining shows the text that has been reversed.

The paragraph embedding level for the first and third examples is 0 (left-to-right direction), and for the second and fourth examples is 1 (right-to-left direction).



Example 1 (embedding level = 0)

Storage:	car means CAR.
Before Reordering:	car means CAR.
Resolved levels:	0000000001110
Reverse level 1:	car means <u>RAC</u> .





Example 2 (embedding level = 1)

Storage:	 car MEANS CAR. 
Before Reordering:	car MEANS CAR.
Resolved levels:	22211111111111
Reverse level 1:	<u>rac</u> MEANS CAR.
Reverse levels 1-2:	<u>.RAC SNAEM car</u>

Example 3 (embedding level = 0)

Storage:	he said "  car MEANS CAR.  .
Before Reordering:	he said "car MEANS CAR."
Resolved levels:	0000000022211111111100
Reverse level 1:	he said " <u>rac</u> MEANS CAR."
Reverse levels 1-2:	he said " <u>RAC SNAEM car</u> ."

Example 4 (embedding level = 1)

Storage:	DID YOU SAY '  he said "  car MEANS CAR.   '?
Before Reordering:	DID YOU SAY 'he said "car MEANS CAR"'?

Resolved levels::	1111111111111222222222444333333333211
Reverse level 4:	DID YOU SAY 'he said " <u>rac</u> MEANS CAR"\'?
Reverse levels 3–4:	DID YOU SAY 'he said " <u>RAC SNAEM car</u> "\'?
Reverse levels 2–4:	DID YOU SAY ' " <u>rac</u> MEANS CAR" dias eh'?
Reverse levels 1–4:	? 'he said "RAC SNAEM car"' YAS UOY DID

L3. Combining marks applied to a right-to-left base character will at this point precede their base character. If the rendering engine expects them to follow the base characters in the final display process, then the ordering of the marks and the base character must be reversed.

Many font designers provide default metrics for combining marks that support rendering by simple overhang. Because of the reordering for right-to-left characters, it is common practice to make the glyphs for most combining characters overhang to the left (thus assuming the characters will be applied to left-to-right base characters) and make the glyphs for combining characters in right-to-left scripts overhang to the right (thus assuming that the characters will be applied to right-to-left base characters). With such fonts, the display ordering of the marks and base glyphs may need to be adjusted when combining marks are applied to “unmatching” base characters. See *Section 5.13, Rendering Nonspacing Marks* of [\[Unicode\]](#), for more information.

L4. A character is depicted by a mirrored glyph if and only if (a) the resolved directionality of that character is R, and (b) the Bidi_Mirrored property value of that character is true.

- *The Bidi_Mirrored property is defined by Section 4.7, Bidi Mirrored—Normative of [\[Unicode\]](#); the property values are specified in [\[UCD\]](#).*
- *This rule can be overridden in certain cases; see [HL6](#).*

For example, U+0028 LEFT PARENTHESIS—which is interpreted in the Unicode Standard as an opening parenthesis—appears as “(” when its resolved level is even, and as the mirrored glyph “)” when its resolved level is odd. Note that for backward compatibility the characters U+FD3E (() ORNATE LEFT PARENTHESIS and U+FD3F ()) ORNATE RIGHT PARENTHESIS are not mirrored.

3.5 Shaping

Cursively connected scripts, such as Arabic or Syriac, require the selection of positional character shapes that depend on adjacent characters (see *Section 8.2, Arabic* of [\[Unicode\]](#)). Shaping is logically applied *after* the Bidirectional Algorithm is used and is limited to characters within the same directional run. Consider the following example string of Arabic characters, which is represented in memory as characters 1, 2, 3, and 4, and where the first two characters are overridden to be LTR. To show both paragraph directions, the next two are embedded, but with the normal RTL direction.

1	2	3	4
ج	ع	ل	م
062C JEEM	0639 AIN	0644 LAM	0645 MEEM
L	L	R	R

One can use embedding codes to achieve this effect in plain text or use markup in HTML, as in the examples below. (The **bold** text would be for the right-to-left paragraph direction.)

- LRM/RLM LRO *JEEM AIN* PDF RLO *LAM MEEM* PDF
- `<p dir="ltr"/>LRO JEEM AIN PDF RLO LAM MEEM PDF</p>`
- `<p dir="ltr"/><bdo dir="ltr">JEEM AIN</bdo>
<bdo dir="rtl">LAM MEEM</bdo></p>`

The resulting shapes will be the following, according to the paragraph direction:

Left–Right Paragraph				Right–Left Paragraph			
1	2	4	3	4	3	1	2
ج	ع	م	ل	م	ل	ج	ع
JEEM-F	AIN-I	MEEM-F	LAM-I	MEEM-F	LAM-I	JEEM-F	AIN-I

3.5.1 Shaping and Line Breaking

The process of breaking a paragraph into one or more lines that fit within particular bounds is outside the scope of the Bidirectional Algorithm. Where character shaping is involved, the width calculations must be based on the shaped glyphs.

Note that the *soft-hyphen*

(SHY) works in cursively connected scripts as it does in other scripts. That is, it indicates a point where the line could be broken in the middle of a word. If the rendering system breaks at that point, the display—including shaping—should be what is appropriate for the given language. For more information on this and other line breaking issues, see Unicode Standard Annex #14, “Line Breaking Properties” [\[UAX14\]](#).

4 Bidirectional Conformance

A process that claims conformance to this specification shall satisfy the the following clauses:

UAX9-C1. *In the absence of a permissible higher-level protocol, a process that renders text shall display all visible representations of characters (excluding format characters) in the order described by Section 3, [Basic Display Algorithm](#), of this annex. In particular, this includes definitions [BD1–BD7](#) and steps [P1–P3](#), [X1–X10](#), [W1–W7](#), [N1–N2](#), [I1–I2](#), and [L1–L4](#).*

- As is the case for all other Unicode algorithms, this is a *logical* description—particular implementations can have more efficient mechanisms as long as they produce the same results. See C18 in *Chapter 3, Conformance of [Unicode]*, and the notes following.
- The Bidirectional Algorithm specifies part of the intrinsic semantics of right-to-left characters and is thus required for conformance to the Unicode Standard where any such characters are displayed.

UAX9-C2. *The only permissible higher-level protocols are those listed in Section 4.3, [Higher-Level Protocols](#). They are [HL1](#), [HL2](#), [HL3](#), [HL4](#), [HL5](#), and*

HL6.

Use of higher-level protocols is discouraged, because it introduces interchange problems and can lead to security problems. For more information, see Unicode Technical Report #36, “Unicode Security Considerations” [UTR36].

4.1 Boundary Neutrals

The goal in marking a format or control character as BN is that it have no effect on the rest of the algorithm. (ZWJ and ZWNJ are exceptions; see X9). Because conformance does not require the precise ordering of format characters with respect to others, implementations can handle them in different ways as long as they preserve the ordering of the other characters.

4.2 Explicit Formatting Codes

As with any Unicode characters, systems do not have to support any particular explicit directional formatting code (although it is not generally useful to include a terminating code without including the initiator). Generally, conforming systems will fall into three classes:

- *No bidirectional formatting.*
This implies that the system does not visually interpret characters from right-to-left scripts.
- *Implicit bidirectionality.*
The implicit Bidirectional Algorithm and the directional marks RLM and LRM are supported.
- *Full bidirectionality.*
The implicit Bidirectional Algorithm, the implicit directional marks, and the explicit directional embedding codes are supported: RLM, LRM, LRE, RLE, LRO, RLO, PDF.

4.3 Higher-Level Protocols

The following clauses are the only permissible ways for systems to apply higher-level protocols to the ordering of bidirectional text. Some of the clauses apply to *segments* of structured text. This refers to the situation where text is interpreted as being structured, whether with explicit markup such as XML or HTML, or internally structured such as in a word processor or spreadsheet. In such a case, a segment is span of text that is distinguished in some way by the structure.

HL1. Override P3, and set the paragraph embedding level explicitly.

- A higher-level protocol may set the paragraph level explicitly and ignore P3. This can be done on the basis of the context, such as on a table cell, paragraph, document, or system level.

HL2. Override W2, and set EN or AN explicitly.

- A higher-level process may reset characters of type EN to AN, or vice versa, and ignore W2. For example, style sheet or markup information can be used within a span of text to override the setting of EN text to be always be AN, or vice versa.

HL3. Emulate directional overrides or embedding codes.

- A higher-level protocol can impose a directional override or embedding on a segment of structured text. The behavior must always be defined by reference to what would happen if the equivalent explicit codes as defined in the algorithm were inserted into the text. For example, a style sheet or markup can set the embedding level on a span of text.

HL4. Apply the Bidirectional Algorithm to segments.

- The Bidirectional Algorithm can be applied independently to one or more segments of structured text. For example, when displaying a document consisting of textual data and visible markup in an editor, a higher-level process can handle syntactic elements in the markup separately from the textual data.

HL5. Provide artificial context.

- Text can be processed by the Bidirectional Algorithm as if it were preceded by a character of a given type and/or followed by a character of a given type. This allows a piece of text that is extracted from a longer sequence of text to behave as it did in the larger context.

HL6. Additional mirroring.

- Characters with a resolved directionality of R that do not have the Bidi_Mirrored property can also be depicted by a mirrored glyph in specialized contexts. Such contexts include, but are not limited to, historic scripts and associated punctuation, private-use characters, and characters in mathematical expressions. (See *Section 6, [Mirroring](#)*.)

Clauses HL1 and HL3 are not logically necessary; they are covered by applications of clauses HL4 and HL5. However, they are included for clarity because they are more common operations.

As an example of the application of HL4, suppose an XML document contains the following fragment. (Note: This is a simplified example for illustration: element names, attribute names, and attribute values could all be involved.)

```
ARABICenglishARABIC<e1 type='ab'>ARABICenglish<e2 type='cd'>english
```

This can be analyzed as being five different segments:

- ARABICenglishARABIC
- <e1 type='ab'>
- ARABICenglish
- <e2 type='cd'>
- english

To make the XML file readable as source text, the display in an editor could order these elements all in a uniform direction (for example, all left-to-right) and apply the Bidirectional

Algorithm to each field separately. It could also choose to order the element names, attribute names, and attribute values uniformly in the same direction (for example, all left-to-right). For final display, the markup could be ignored, allowing all of the text (segments a, c, and e) to be reordered together.

When text using a higher-level protocol is to be converted to Unicode plain text, for consistent appearance formatting codes should be inserted to ensure that the order matches that of the higher-level protocol.

5 Implementation Notes

5.1 Reference Code

There are two versions of BIDI reference code available. Both have been tested to produce identical results. One version is written in Java, and the other is written in C++. The Java version is designed to closely follow the steps of the algorithm as described below. The C++ code is designed to show one of the optimization methods that can be applied to the algorithm, using a state table for one phase.

One of the most effective optimizations is to first test for right-to-left characters and not invoke the Bidirectional Algorithm unless they are present.

There are two directories containing source code for reference implementations: [BidiReferenceJava](#) and [BidiReferenceCpp](#) [Code9]. Implementers are encouraged to use this resource to test their implementations.

5.2 Retaining Format Codes

Some implementations may wish to retain the format codes when running the algorithm. The following provides a summary of how this may be done. Note that this summary is an informative implementation guideline; it should provide the same results as the explicit algorithm above, but in case of any deviation the explicit algorithm is the normative statement for conformance.

- In rule X9, instead of removing the format codes, assign the embedding level to each embedding character, and turn it into BN.
- In rule X10, assign L or R to the last of a sequence of adjacent BNs according to the *eor* / *sor*, and set the level to the higher of the two levels.
- In rule W1, search backward from each NSM to the first character in the level run whose type is not BN, and set the NSM to its type. If the NSM is the first non-BN character, it will get the type of *sor*.
- In rule W4, scan past BN types that are adjacent to ES or CS.
- In rule W5, change all appropriate sequences of ET and BN, not just ET.
- In rule W6, change all BN types adjacent to ET, ES, or CS to ON as well.
- In rule W7, scan past BN.
- In rules N1 and N2, treat BNs adjoining neutrals same as those neutrals.
- In rules I1 and I2, ignore BN.
- In rule L1, include format codes and BN together with whitespace characters in the sequences whose level gets reset before a separator or line break. Resolve any LRE, RLE, LRO, RLO, PDF, or BN to the level of the preceding character if there is one, and otherwise to the base level.

Implementations that display visible representations of format characters will want to adjust this process to position the format characters optimally for editing.

5.3 Joiners

As described under [X9](#), the *zero width joiner* and *non-joiner* affect the shaping of the adjacent characters—those that are adjacent in the original backing-store order—even though those characters may end up being rearranged to be non-adjacent by the Bidirectional Algorithm. To determine the joining behavior of a particular character after applying the Bidirectional Algorithm, there are two main strategies:

- When shaping, an implementation can refer back to the original backing store to see if there were adjacent ZWNJ or ZWJ characters.
- Alternatively, the implementation can replace ZWJ and ZWNJ by an out-of-band character property associated with those adjacent characters, so that the information does not interfere with the Bidirectional Algorithm and the information is preserved across rearrangement of those characters. Once the Bidirectional Algorithm has been applied, that out-of-band information can then be used for proper shaping.

5.4 Vertical Text

In the case of vertical line orientation, the Bidirectional Algorithm is still used to determine the levels of the text. However, these levels are not used to reorder the text, because the characters are usually ordered uniformly from top to bottom. Instead, the levels are used to determine the rotation of the text. Sometimes vertical lines follow a vertical baseline in which each character is oriented as normal (with no rotation), with characters ordered from top to bottom whether they are Hebrew, numbers, or Latin. When setting text using the Arabic script in vertical lines, it is more common to employ a horizontal baseline that is rotated by 90° counterclockwise so that the characters are ordered from top to bottom. Latin text and numbers may be rotated 90° clockwise so that the characters are also ordered from top to bottom.

The Bidirectional Algorithm is used when some characters are ordered from bottom to top. For example, this happens with a mixture of Arabic and Latin glyphs when all the glyphs are rotated uniformly 90° clockwise. The Unicode Standard does not specify whether text is presented horizontally or vertically, or whether text is rotated. That is left up to higher-level protocols.

5.5 Usage

Because of the implicit character types and the heuristics for resolving neutral and numeric directional behavior, the implicit bidirectional ordering will generally produce the correct display without any further work. However, problematic cases may occur when a right-to-left paragraph begins with left-to-right characters, or there are nested segments of different-direction text, or there are weak characters on directional boundaries. In these cases, embeddings or directional marks may be required to get the right display. Part numbers may also require directional overrides.

The most common problematic case is that of neutrals on the boundary of an embedded language. This can be addressed by setting the level of the embedded text correctly. For example, with all the text at level 0 the following occurs:

Memory: he said "I NEED WATER!", and expired.

Display: he said "RETAW DEEN I!", and expired.

If the exclamation mark is to be part of the Arabic quotation, then the user can select the text *I NEED WATER!* and explicitly mark it as embedded Arabic, which produces the following result:

Memory: he said "<RLE>I NEED WATER!<PDF>", and expired.

Display: he said "!RETAW DEEN I", and expired.

However, a simpler and better

method of doing this is to place a right directional mark (RLM) after the exclamation mark. Because the exclamation mark is now not on a directional boundary, this produces the correct result.

Memory: he said "I NEED WATER!<RLM>", and expired.

Display: he said "!RETAW DEEN I", and expired.

This latter approach is preferred because it does not make use of the stateful format codes, which can easily get out of sync if not fully supported by editors and other string manipulation. The stateful format codes are generally needed only for more complex (and rare) cases such as double embeddings, as in the following:

Memory: DID YOU SAY '<LRE>he said "I NEED WATER!<RLM>", and expired.<PDF>'?

Display: ?'he said "!RETAW DEEN I", and expired.' YAS UOY DID

5.6 Separating Punctuation Marks

A common problem case is where the text really represents a sequence of items with separating punctuation marks, often programmatically concatenated. These separators are often strings of neutral characters. For example, a web page might have the following at the bottom:

advertising programs - business solutions - privacy policy - help - about

This might be built up on the server by concatenating a variable number of strings with " - " as a separator, for example. If all of the text is translated into Arabic or Hebrew and the overall page direction is set to be RTL, then the right result occurs, such as the following:

TUOBA - PLEH - YCILOP YCAVIRP - SNOITULOS SSENISUB - SMARGORP
GNISITREVDA

However, suppose that in the translation, there remain some LTR characters. This is not uncommon for company names, product names, technical terms, and so on. If one of the separators is bounded on both sides by LTR characters, then the result will be badly jumbled. For example, suppose that "programs" in the first term and "business" in the second were left in English. Then the result would be:

TUOBA - PLEH - YCILOP YCAVIRP - SNOITULOS programs - business GNISITREVDA

The result is a jumble, with the apparent first term being "advertising business" and the second being "programs solutions". The simplest solution for this problem is to include an RLM character in each separator string. That will cause each separator to adopt a right-to-left direction, and produce the correct output:

TUOBA - PLEH - YCILOP YCAVIRP - SNOITULOS business - programs GNISITREVDA

The stateful controls (LRE, RLE, and PDF) can be used to achieve the same effect; web pages would use spans with the attributes *dir="ltr"* or *dir="rtl"*. Each separate field would be embedded, excluding the separators. In general, LRM and RLM are preferred to the stateful approach because their effects are more local in scope, and are more robust than the *dir* attributes when text is copied. (Ideally programs would convert *dir* attributes to the

corresponding stateful controls when converting to plain text, but that is not generally supported.)

Migrating from 2.0 to 3.0

In the Unicode Character Database for [Unicode3.0](#), new bidirectional character types were introduced to make the body of the Bidirectional Algorithm depend only on the types of characters, and not on the character values. The changes from the 2.0 bidirectional types are listed in *Table 6*.

Table 6. New Bidirectional Types in Unicode 3.0

Characters	New Bidirectional Type
All characters with General_Category Me, Mn	NSM
All characters of type R in the Arabic ranges (0600..06FF, FB50..FDFF, FE70..FEFE) (Letters in the Thaana and Syriac ranges also have this value.)	AL
The explicit embedding characters: LRO, RLO, LRE, RLE, PDF	LRO, RLO, LRE, RLE, PDF, respectively
Formatting characters and controls (General_Category Cf and Cc) that were of bidirectional type ON	BN
Zero Width Space	BN

Implementations that use older property tables can adjust to the modifications in the Bidirectional Algorithm by algorithmically remapping the characters in *Table 6* to the new types.

6 Mirroring

The mirrored property is important to ensure that the correct character codes are used for the desired semantic. This is of particular importance where the name of a character does not indicate the intended semantic, such as with U+0028 “(” LEFT PARENTHESIS. While the name indicates that it is a left parenthesis, the character really expresses an *open parenthesis*—the *leading* character in a parenthetical phrase, not the trailing one.

Some of the characters that do not have the Bidi_Mirrored property may be rendered with mirrored glyphs, according to a higher level protocol that adds mirroring: see *Section 4.3, Higher-Level Protocols*, especially [HL6](#). Except in such cases, mirroring must be done according to rule [L4](#), to ensure that the correct character code is used to express the intended semantic of the character, and to avoid interoperability and security problems.

Implementing rule [L4](#) calls for mirrored glyphs. These glyphs may not be exact *graphical* mirror images. For example, clearly an italic parenthesis is not an exact mirror image of another—“(” is not the mirror image of “)”. Instead, mirror glyphs are those acceptable as mirrors within the normal parameters of the font in which they are represented.

In implementation, sometimes pairs of characters are acceptable mirrors for one another—for example, U+0028 “(” LEFT PARENTHESIS and U+0029 “)” RIGHT PARENTHESIS or U+22E0 “⋈” DOES NOT PRECEDE OR EQUAL and U+22E1 “⋉” DOES NOT SUCCEED OR EQUAL. Other characters such as U+2231 “∫” CLOCKWISE INTEGRAL

do not have corresponding characters that can be used for acceptable mirrors. The informative Bidi Mirroring data file [\[Data9\]](#), lists the paired characters with acceptable mirror glyphs. A comment in the file indicates where the pairs are “best fit”: they should be acceptable in rendering, although ideally the mirrored glyphs may have somewhat different shapes.

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References

For references for this annex, see Unicode Standard Annex #41, “[Common References for Unicode Standard Annexes](#).”

Modifications

The following summarizes modifications from previous revisions of this annex.

Revision 18

- Clarified [BD6](#).
- Made some examples more explicit.
- Added the common problem case of separators in 5.5.
- Added notes on security concerns with RLO and LRO, and the use of dir="ltr" or "rtl" with web pages.
- Fixed example under [N2](#).
- **Draft 4**
 - Fixed example in Section 3.3.4 [Resolving Neutral Types](#).
 - Made last part of Section 5.5 [Usage](#) into 5.6 [Separating Punctuation Marks](#), changed the term "Separators" and added note on stateful controls.

Revision 17

- This revision incorporates successive changes. The latest changes were on 2006-02-24.
- Modified [L4](#) and [HL6](#), in conjunction with proposed property change to Bidi_Mirrored ([PRI #80](#)).
- Added note on U+ FD3E (()) ORNATE LEFT PARENTHESIS and U+ FD3F (()) ORNATE RIGHT PARENTHESIS.
- Used new format for conformance clause numbering.
- Added caution on use of higher-level protocols, after [UAX9-C2](#).
- Some wording changes in 6 [Mirroring](#), for consistency with new L4 and HL6.
- Moved text to [Shaping and line breaking](#), and added note on SHY in 3.4 there.

- Removed two notes indicating that the conformance clauses override clause C13 of Unicode 4.0.
- Changed some references to Unicode4.0

Revision 16 being a proposed update, only changes between versions 17 and 15 are noted here.

Revision 15:

- Minor editing
- Fixed section Number for Mirroring
- Changed “Tracking Number” to Revision
- Added note on U+0CBF KANNADA VOWEL SIGN I
- Added note after N1, and clarified example after N2.
- Fixed references to sections of the Unicode Standard

Revision 14:

- Aliased directional run and level run
- Pointed to DerivedBidiClass.txt for unassigned character assignments.

Revision 13:

- [4. Bidirectional Conformance](#): added explicit clauses.
- [4.3. Higher-Level Protocols](#):
 - Added clarifying text, and renumbered options.
 - Removed option regarding number shaping (because it was irrelevant to bidirectional ordering).
 - Broadened the ability to override on the basis of context, and clarified number handling.
 - Made clear that bidi could be applied to segments
- [1. Introduction](#): added note that the changes in [4. Bidirectional Conformance](#) override clause C13 of Unicode 4.0 [[Unicode](#)], and tighten the conformance requirements from what they had been previously.
- Minor editing for clarification.

Revision 11:

- Updated for Unicode 4.0.
- Added note on [canonical equivalence](#)
- Added [Joiners](#) section on ZWJ and ZWNJ
- Clarified [L2](#) and examples following.
- Added a section on the interaction of [shaping](#) and bidirectional reordering.
- Moved lists for unassigned characters into UCD.html (also now explicit in DerivedBidiClass.txt)
- Updated references for Newline Guidelines (because the UAX is incorporated into the 4.0 book)
- The first two sections were rearranged, with [Reference Code](#) going into [Implementation Notes](#), and [Mirroring](#) in its own section at the end.
 - This is *not* highlighted in the proposed text.
- Sections were renumbered and the table of contents is more detailed.

- This is *not* highlighted in the proposed text.
- Misc editing.

Revision 10:

- Updated for Unicode 3.2.
- Updated UAX boilerplate in the status section.

Revision 9:

- Clarified the language of [P2](#)
- Corrected the implementation note on “Retaining Format Codes” in [Implementation Notes](#)
- Minor editing

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